

The hour glass method for encryption and compressing any encrypted text into a single character, using the fundamentals of power arithmetic function to reduce space complexity

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Abstract: Encryption is the science of concealing data, and because researchers seek to find the best way to hide data from unauthorized persons to access it, and due to competition in this field, it became necessary to find methods that are difficult for attackers to break the cipher text and know its content, in this research paper a new method was proposed that we called it The "hour glass or sand clock method", which mainly depends on reducing the size of any message, regardless of its size, to become eight-bit size only, in addition to hiding all characters of the message with only one character using the principle of arithmetic power function, and it is the first method that relies in its work on the principle of compressing a complete message to become

Only one character. Which poses a real challenge in front of all methods of breaking the cipher text as it will deal with only one character. By using different cryptanalysis methods, the proposed method's strength and effectiveness have been proven in countering all types of attacks.

Keywords: sand clock, hour glass, power function, single character, compression

1. Introduction

Hour glass or sand clock

The hourglass is a clock that consists of two glass containers, an upper and lower connected, open in the middle and placed in the upper part of sand, and when the sand reaches the hole between them, the passage of sand becomes one atom after another until it accumulates a second time in the lower part figure(1).

Figure1: hour glass or sand clock



The encryption and compression in the proposed method behave as the sand in the upper container, and the decryption process and decompression behave as the grouping of the sand atoms in the lower container.

Power function: The power function in mathematics can be defined as $f(x) = a^x$ where $a, x \in R$, there are two types of power function which is the odd function and even function as shown in the figures (2) & (3)

Figure2:evenfunctiontype

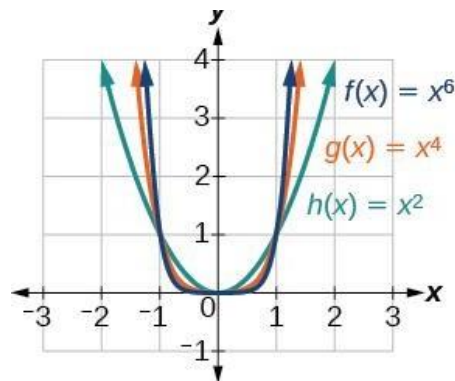
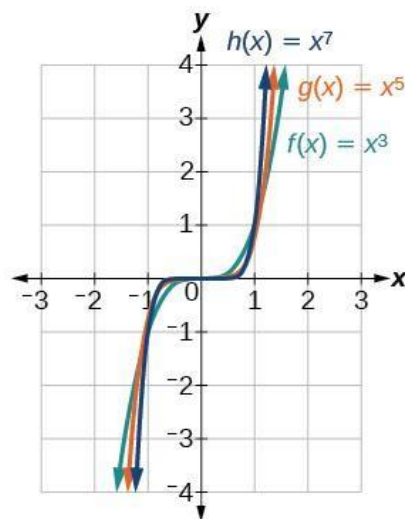


Figure3:oddfunctiontype



the power function behave mathematically according to the value of the power ,which means when the power of the basis start from 0 to $n \in R$, the value of the power function will increase accordingly, the power function represent the product of the basis by itself times number according to the power of the basis $a^n = a \times a \times \dots n \text{ times}$

also there is exponential function and power function for the basis 10, as another types of power function (Zaka, A., Akhter, A. S., & Jabeen, R. (2020))

Cryptography

Cryptography is the science that deals with the cryptographic systems.

Cryptanalysis is the technique of breaking the cryptographic systems.

Cryptography is one of the fields of computer science & mathematics that focuses on all techniques that made a secure communication between two persons (Alice & Bob)

The four principles of cryptography are:

Confidentiality: Defines a set of the basis that limits access and/or adds restriction on the certain information.

Data Integrity: deals with the consistency and the accuracy of the data along its life-cycle.

Authentication : it confirms the truth of the attribute of a datum which is cleared to be “True” through using some entity.

Non-Repudiation : To ensure the inability of the author of the statement resp. which is the piece of the information to be denied. today there are two different schemes , which is the symmetric schemes , based on both Alice and Bob need to obtain the same key to encrypt and decrypt their communication .So, they need to share the key initially. On the other side, Diffie and Hellman's was invited the key exchange idea, so the concept of asymmetric schemes which is Alice & Bob have two keys, the private and public key. The public key can be shared with any person or any one , this public key can be used by Bob to encrypt the secret message and send it for Alice. In the other hand only Alice can use the corresponding private key to decrypt the cipher text from Bob. (**Barakat, M., Eder, C., & Hanke, T. (2018)**)

2. Significance Of The Study

Most of the secret messages can be breaking through the difference types of the cryptanalysis methods , and all of these messages has the same number of the characters , that make the guess operation very simple by comparing each character with the equivalent character in the secret message , but in this new method all message characters will be gathering into one character , this property make the knowing of the original message very difficult, so all the cryptanalysis will fail to guess the correct number of the character and doesn't decipher the message.

3. Review Of Related Studies

Kumari, M., & Tanti, J. (2020) proposed a Public Key using the block Cryptography matrices with generalized Fibonacci sequence. That show first the multiplicative commutativity of the generalized matrices that constructed through using the generalized Fibonacci sequences, then they developed an cryptographical scheme , they also discussed the efficiency & the strength of the proposed scheme in the context of the block matrices.

Abed, B. N., et.al (2020) proposed a new cryptosystem through using the 3rd order equation "cubic equation" and use Cardano's method , the purpose is to add more secrecy and complexity to the proposed cryptographic algorithm, while there is four different keys and different formulas of the equation . **Noaman, S. A., et . al (2020)** , presents the method of data encryption / decryption. The proposed method used the Taylor series through choosing the constant and Taylor formula as two secret keys. the first key substitutes the plain text in the Taylor. Where in Decoding phase compute the Taylor inversion. many attacks, are used in order to evaluate the results of the algorithm.

Najim al-din, B., & Shaban, S. A. (2017) proposed the cryptosystem to encrypt Arabic text through the principle of integration to produce better security as well as increasing the complexity of predicting the secret keys and know the true plaintext. the results show that the proposed cryptosystem was inevitable to cryptanalysis process.

Abed, B. N. A. D., & Noaman, S. A. (2019) presents a new method to develop the techniques of encryption , through the McLaurin series considering as a new cryptosystem , by using different cryptanalysis techniques with different decoding tools , the proposed algorithm was inevitable against all these different attacks , in addition it proved as a one way function . **Al-din, B. N., et.al. (2020)** proposes an algorithm known as wolf algorithm, through classifying characters of the plain text into many groups and then exchange the keys between all these groups in the wolf communication process to construct an authentication secret keys between the groups , many different cryptanalysis techniques used to evaluate the proposed new algorithm.

4. Objectives Of The Study

- to find a new method that can hide the information of the secret message
- to compress the message into one character
- to decipher the message using the send one character instead of complete message
- to make the cryptosystem very complex and difficult for breaking

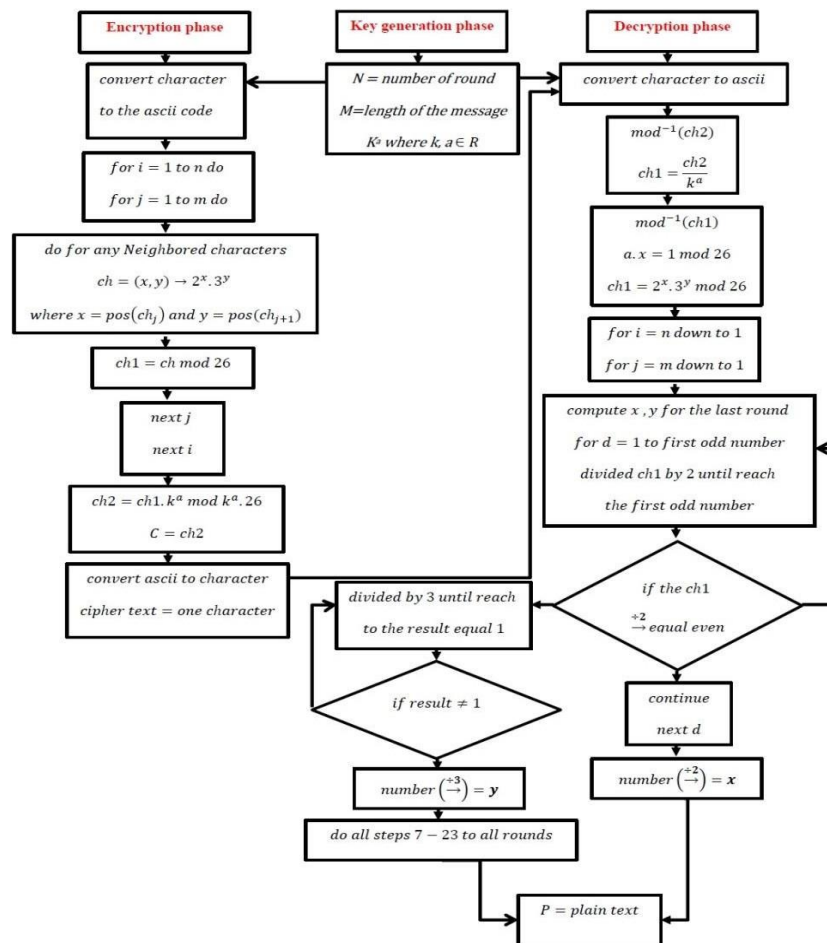
5.HypothesesOfTheStudy

- thereisnocryptosystemhidethemessageinverysecure wayforever
- thereisnocryptosystemhidethemessageinonecharacter
- thereisnocryptosystemcompresses the secret messagein order todecriesthememoryusage

6.Methodology

In this section, three algorithms are presented that represent the three aspects of the cryptographic and compression system, which are the key generation phase, the encryption phase, and the decryption phase. These three aspects are illustrated in the flowchart of the proposed method. In addition to listing all these algorithms in detail, and making an implementation of the proposed method in order to list the complete details of the method practically and mathematically to add reliability to the proposed method. And as shown in the practical example in this section and the results achieved theoretically and practically to reach the encrypted text, which is represented by only one character, which is at the same time compressing the size of a message to a much smaller size, which represents only one byte, which will reduce the storage space consumed in addition to increasing security. And the biography of the transmitted data via the Internet.

Figure4:flowchartoftheproposedmethod



Keygeneration

n is the number of the steps that execute the algorithm
m is the length of the message
k^a is the third secret key where $k, a \in \mathbb{R}$

Encryption

```

convert character to the ascii code for
i = 1 to n do
  for j = 1 to m do
    do for any Neighbored characters
       $ch = (x, y) \rightarrow 2^x \cdot 3^y$  where  $x = \text{pos}(ch_j)$  and  $y = \text{pos}(ch_{j+1})$ 
       $ch1 = ch \bmod 26$ 
      next j
    next i
     $ch2 = ch1 \cdot k^a \bmod k^a \cdot 26$ 
     $C = ch2$ 
  convert ascii to character
  cipher text = one character

```

Decryption

```

convert character to ascii
 $ch2 = \frac{\text{mod}^{-1}(ch2)}{k^a}$ 
 $ch1 = \text{mod}^{-1}(ch1)$ 
 $a \cdot x = 1 \bmod 26$ 
 $ch1 = 2^x \cdot 3^y \bmod 26$ 
for i = n down to 1
  for j = m down to 1
    compute x, y for the last round for d =
      1 to first odd number
      divided ch1 by 2 until reach the first odd number if the c
      h1 divided by 2 equal even number
    continue
  next d
  else stop
  the number of the divided by 2 is the x value
  if the result of divided by 2 equal odd number then for
    l = 1 to the result equal 1
    divided by 3 until reach to the result equal 1 if the r
    esult not equal 1 continue
  else stop
  the number of the divided by 3 is the y value
  next l
do all steps 7 – 23 to all rounds
next m

```

nextn

P=plaintext

7.Implementation

Inthisalgorithmwewilladoptnewalphanumeric values that differ from the ASCII code, as shown in the table below

Figure5:thealphabettable

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26

Key generation

n=4 where n is the number of rounds for encryption

m=11 where m is the length of the message

$k^a=2^2$ where $a=2$, $k=2$ the third secret key

Encryption

P="GOODMORNING" //where P is the plaintext

P= 7 15 15 4 13 15 18 14 9 14 7 // Ascii

code in the first round

$(7,15) \rightarrow 2^7 \cdot 3^{15} = 1836660096 \bmod 26 = 24$

in the same method $(15,4)=24$, $(13,15)=2$, $(18,14)=4$, $(9,14)=6$, and 7

the result of the first round is: 24 24 24 6 7 in

the second round

$(24,24)=14$, $(2,4)=12$, $(6,7)=10$

the result of the second round is: 14 12 10 in the

third round

$(14,12)=4$, 10

the result in the third round is:

4 10 in the fourth and final round

$(4,10)=22=V$

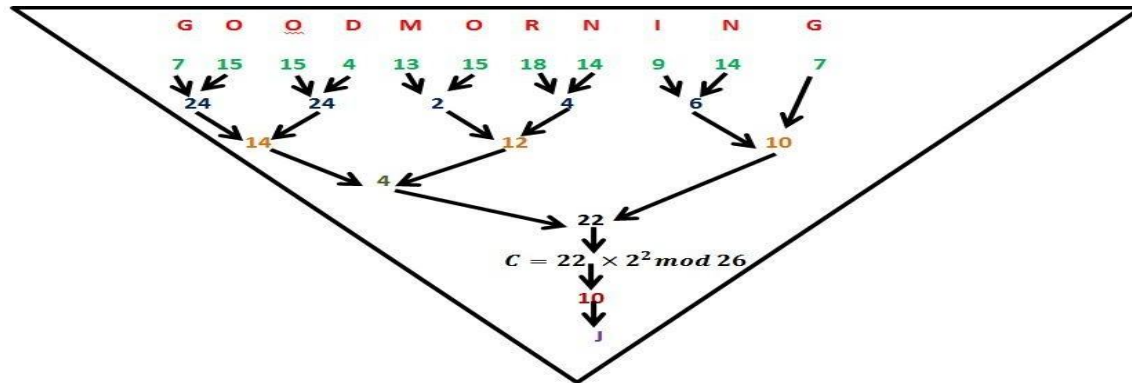
$C=2^2 \cdot 22 \bmod 26 = 10$

$C(10)=J$

ciphertext=*J*

thefollowingchartrepresent“thehourglass”methodforencryptionandcompression

Figure6:encryptionprocesses



Decryption

J=10

$\text{mod}^{-1}(10)=22$

$\text{mod}^{-1}(22)=944784$

inthelastroundandcomputex,yvalues

$944784 \xrightarrow{\div 2} 472392 \xrightarrow{\div 2} 236196 \xrightarrow{\div 2} 118098 \xrightarrow{\div 2} 59049$ is odd stop then x value
 $=4$ (number of divided by 2 \rightarrow)

$59049 \xrightarrow{\div 3} 19683 \xrightarrow{\div 3} 6561 \xrightarrow{\div 3} 2187 \xrightarrow{\div 3} 729 \xrightarrow{\div 3} 243 \xrightarrow{\div 3} 81 \xrightarrow{\div 3} 27 \xrightarrow{\div 3} 9 \xrightarrow{\div 3} 3 \xrightarrow{\div 3} 1$ the result is one stop then y value
 $=10$ (number divide by 3 \rightarrow)

theresultfromthereversefirstroundis:410

theresultfromthereversesecoundroundis:141210

theresultfromthereversethirdroundis:24242467

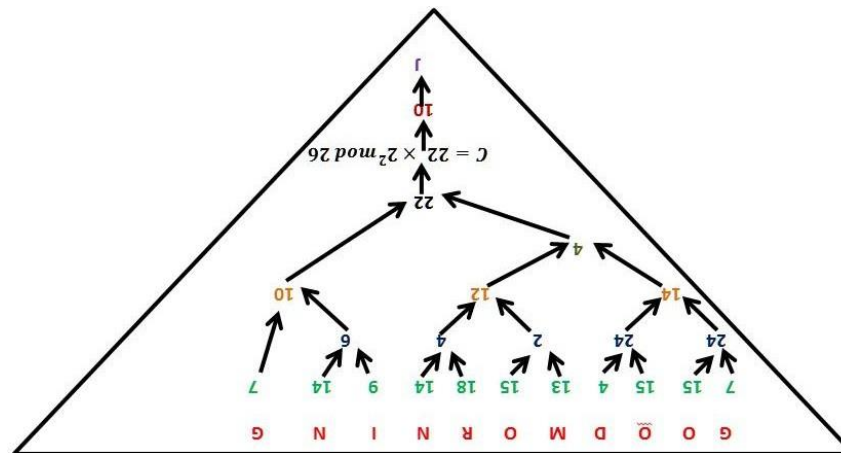
theresultfromthereversefourthroundis:715154131518149147

throughconvertasciitothecharacterwillbe:GOODMORNING

P=GOODMORNING

Thechartbelowrepresentthederyptionanddecompressionprocesses

Figure7:decryptionprocesses



8.Resultanddiscussion

In order to prove that the proposed method is efficient and effective, a set of crypt analysis methods havebeenadoptedthat willprovewhetheritisabletoguessandpredicttheexplicittextand thekeyusedintheencryptionprocess,whichare asfollows:

When applying frequency analysis to the encrypted text consisting of only one character, this method provedits inability to break the encrypted text as it showed that the cipher text is treated as only one character, which is "J" inthis example, so it is difficult for this method to know the real number of characters, and what are the contents of thismessageasinthe figures(8), (9)

Figure8:frequencyanalysis

Frequencies:

Single letters:

J	A	B	C	D	E	F	G	H	I	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
100%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%

Bigrams:

AA	AB	AC	AD	AE	AF	AG	AH	AI	AJ	AK	AL	AM	...
0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	0%	...

Guesses (clear guesses):

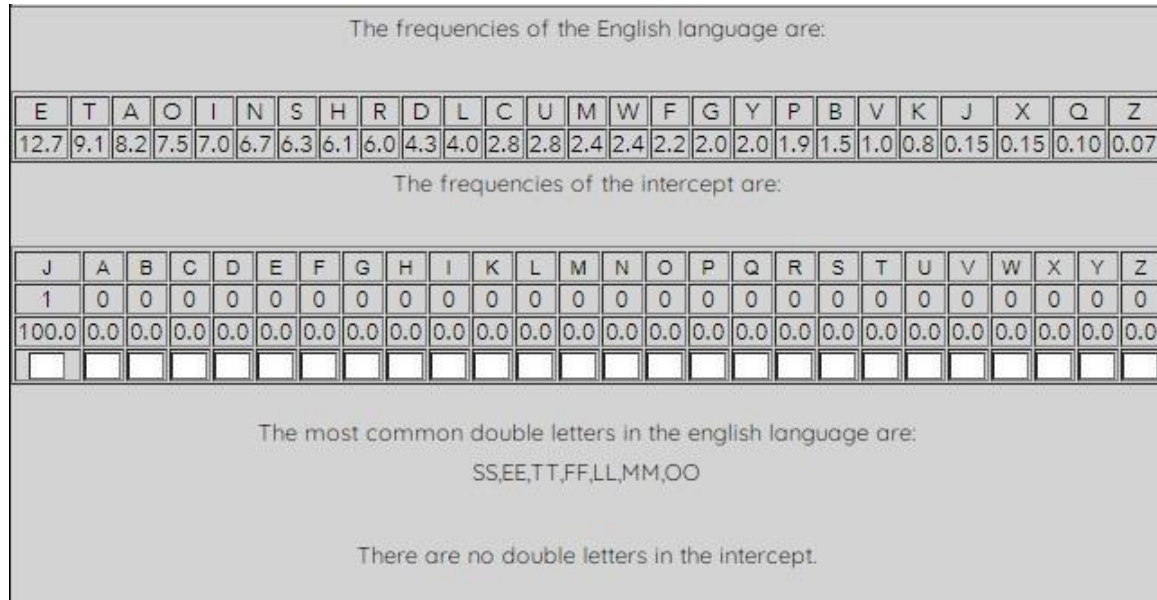
A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Figure9:frequencycomparission

E	T	A	O	I	N	S	H	R	D	L	C	U	M	W	F	G	Y	P	B	V	K	J	X	Q	Z
12.7	9.1	8.2	7.5	7.0	6.7	6.3	6.1	6.0	4.3	4.0	2.8	2.8	2.4	2.4	2.2	2.0	2.0	1.9	1.5	1.0	0.8	0.15	0.15	0.10	0.07
The frequencies of the intercept are:																									
J	A	B	C	D	E	F	G	H	I	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
100.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0

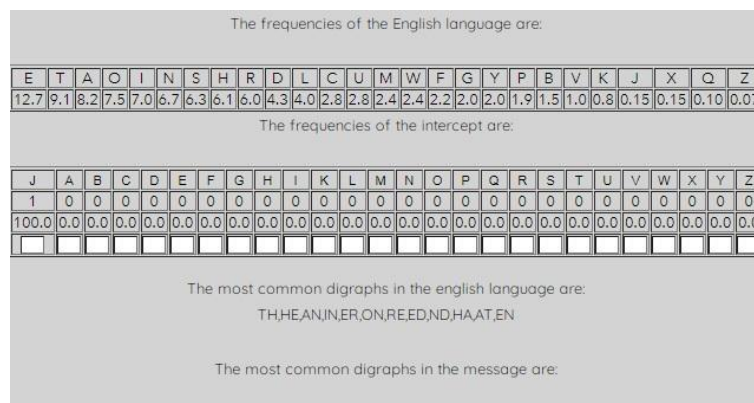
In addition to the inability to determine the most common double letters as is recognized in the English letters, as the results showed that it was not possible to specify any double letter as shown in Figure(10)

Figure10:the most common double letters



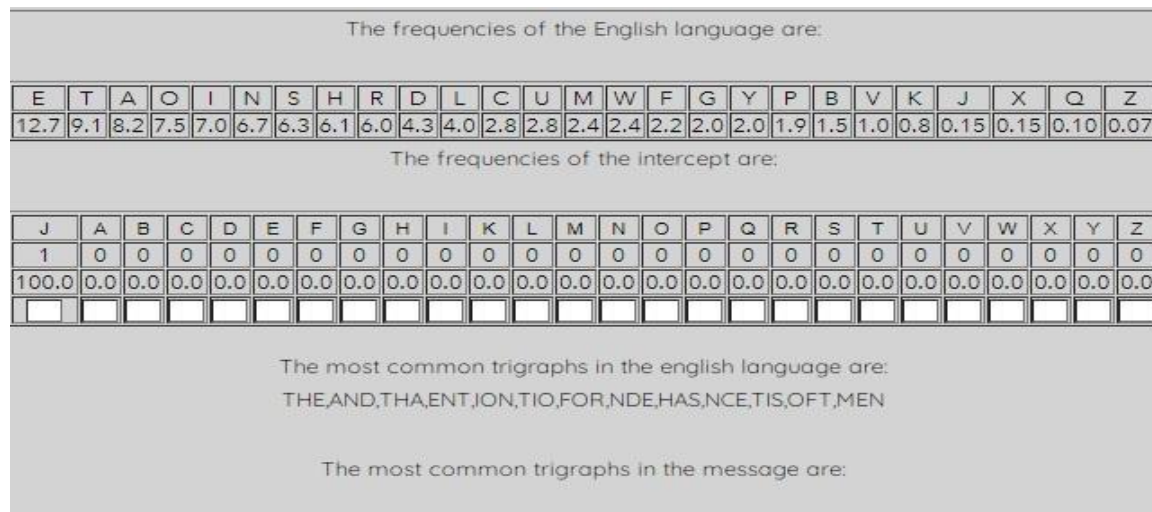
Likewise, the inability to determine the most common digraphs characters as it is known in English letters or what is known conventionally, as the results showed that it was not possible to specify any digraphs as shown in Figure(11)

Figure11:the most common digraphs



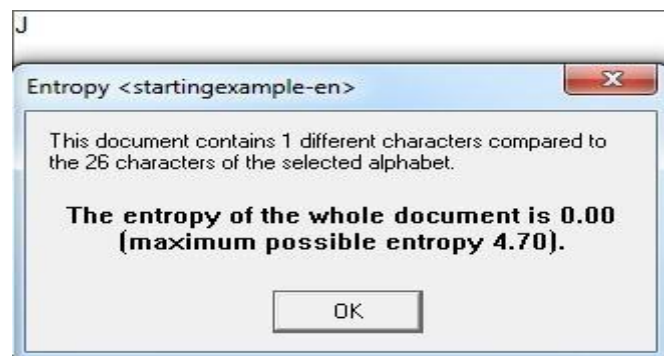
the same thing for the three letters..as it proved that the most common three letters cannot be identified as it is known in the English letters or what is known conventionally trigraphs , as the results showed that it was not possible to specify any three letters as shown in Figure (12)

Figure12:themostrcommontrigraphs



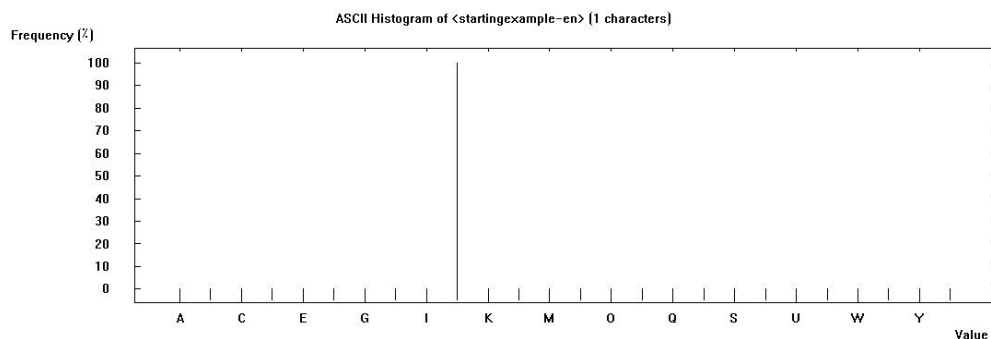
Also,theentropyoftheencryptedtextshowedtheprobabilityofappearingfortheletterwithavalueof0.0because them essagcontainsonlyonelettercomparedto 26letters, asinthe figure(13)

Figure13:theentropy



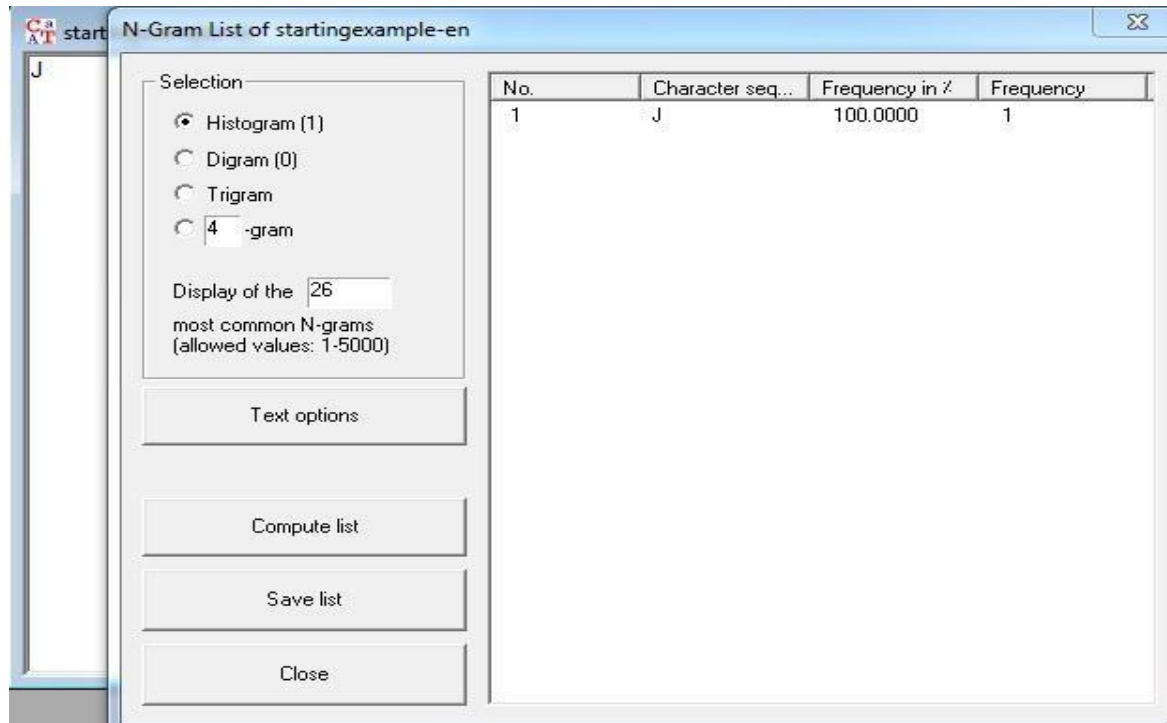
And the analysis of the histogram revealed the presence of the frequency ratio, which is also due to thepresenceofonlyoneletter intheletter,whichmakestheprocessofdeterminingtheproportionsoftheotherlettersdifficultandasshowninthe histograminFigure (14)

Figure14:thehistogram



Likewise, the N-Gram analysis showed that the frequency of the letter J is 100%, and this indicates that it deals with the encoded text as only one letter, and this is not identical to reality as the number of letters in this test is 11 letters and not one letter, so this analysis failed. Also, in knowing the explicit text as in the figure (15)

Figure15:N-Gram



From all of the above mentioned tests that were done on the encrypted text, all the methods dealt with the encrypted text on the basis of a single character, and this proves that any other method will deal in the same way with the encrypted text, which makes the process of analyzing and decrypting the cipher text very complicated or seem impossible. For this method.

The space complexity of any algorithm is the total amount of computer memory that required by an algorithm in order to complete its execution. If the amount of the space that required by the algorithm increasing with the increase of the input value, as a result the space complexity is "Linear Space Complexity". And if the algorithm that requires a fixed amount of space, for all the input values, the space complexity is a "Constant Space Complexity". (Queiroz, S., Silva, W., Vilela, J. P., & Monteiro, E. (2020), Can, T., Krishnamurthy, K., & Schwab, D. J. (2020, August))

The time complexity of any algorithm is said to be the total amount of the time that required by the algorithm in order to complete program execution. And if the program requires the fixed amount of the time for all the input values is lead to say that the time complexity is "Constant Time Complexity". Where if the amount of time that required by any algorithm increasing with the increase of the input value, the time complexity is "Linear Time Complexity". So the time complexity for the key generation algorithm is $O(1)$ and space complexity is $O(1)$ which the best and the time complexity for encryption and decryption algorithms is $O(n)$, and space complexity is $O(n)$.

We also compared the proposed method with the methods in the literature review in terms of reducing the size of the message and in terms of the keys, as in Table No. (1)

Table1:compression table

method	Number of keys	Reduce the size	compressed	efficiency	Number of Ciphertext characters
Kumari, M., & Tanti, J. (2020)	2	No	No	Yes	Equal plain
Abed, B. N., et. al (2020)	3	No	No	Yes	Equal plain
Noaman, S. A., et. al (2020)	2	No	No	Yes	Equal plain
Najim al-din, B., & Shaban, S. A. (2017).	2	No	No	Yes	Equal plain
Abed, B. N. A. D., & Noaman, S. A. (2019)	3	No	No	Yes	Equal plain
Al-din, B. N., Manasrah, A. M., & Noaman, S. A. (2020)	2	No	No	Yes	Equal plain
Proposed method	3	Yes	Yes	Yes	One character

9. Conclusion

In this research paper, a method called the hourglass method was proposed to build a highly efficient encryption system based on the principle of data compression, which is also an effective method of pressure that achieves the least possible space, which saves large storage space and reduces the space complexity very significantly, in addition to security and confidentiality. The high level of this system, which will preserve the data against various types of attacks, thus providing a safe way to transfer data. The great contribution made by this method in the field of data security science and its compression will open great horizons for researchers, and many methods of cryptanalysis have been applied on the proposed system and the results proved the strength and effectiveness of this method and the inability of the cryptanalysis methods to break it, in the future work we can apply this method to compress and steganography's the images and videos and other resources.

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